

What is a cell cycle checkpoint: The TotemBioNet answer

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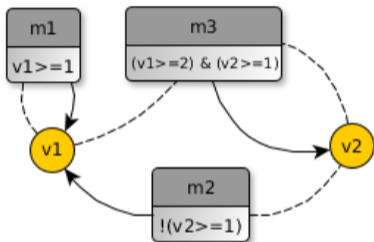
Région
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TotemBioNet and the qualitative modeling framework

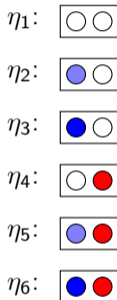
- TotemBioNet automates parameters' identification for René Thomas' discrete modeling framework
- It combines two formal methods: weakest precondition for Hoare logic and model checking for temporal logic
- It computes the *exhaustive* set of Thomas' parameterizations verifying a set of biological properties

René Thomas' syntax: multivaluated regulatory network

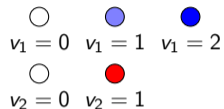
Regulatory graph:



States of the system:



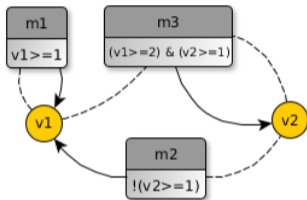
States of variables:



- \mathbf{m}_1 $[v_1 \geq 1] \rightarrow v_1$: v_1 activates itself
- \mathbf{m}_2 $[\neg(v_2 \geq 1)] \rightarrow v_1$: v_2 inhibits v_1
- \mathbf{m}_3 $[(v_1 \geq 2) \wedge (v_2 \geq 1)] \rightarrow v_2$: an activating dimer of v_2

René Thomas' Semantic: asynchronous dynamic

Regulatory graph:



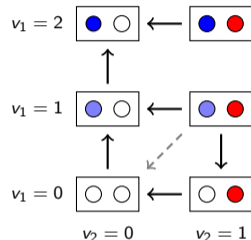
Set of resources:

- ω_{v_1} : m_1, m_2
- ω_{v_2} : m_3

K parameters:

$$\begin{aligned}
 K_{v_1, \emptyset} &= 0 \\
 K_{v_1, m_1} &= 0 \\
 K_{v_1, m_2} &= 1 \\
 K_{v_1, m_1, m_2} &= 2 \\
 K_{v_2, \emptyset} &= 0 \\
 K_{v_2, m_3} &= 1
 \end{aligned}$$

Asynchronous transition graph:

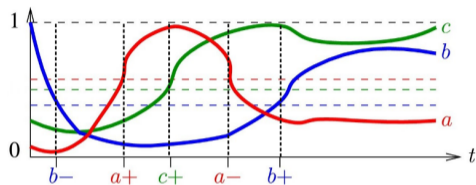


- System dynamics depends on K parameters of the form K_{v, ω_v} where ω_v is a resource of v .
- From a regulatory graph, number of parameterizations : $\prod_v (d^+(v) + 1)^{2^{d^-(v)}}$ where $d^+(v)$ and $d^-(v)$ are resp. the outdegree and indegree of v .

The genetically modified Hoare logic

Hoare triple noted H: {Pre} Path {Post}

- Precondition: $a=0, b=1, c=0$
- Path: $b-; a+; c+; a-; b+$
- Postconditions: $a=0, b=1, c=1$



Normalised expression profiles from a biological experiment

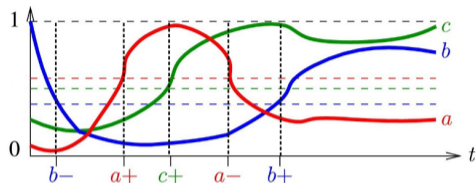
$$H_{ex} : \{a = 0, b = 1, c = 0\} \quad b- ; a+ ; c+ ; a- ; b+ \quad \{a = 0, b = 1, c = 1\}$$

↓
Postcondition noted Q

The genetically modified Hoare logic

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- Precondition: $a=0, b=1, c=0$
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Normalised expression profiles from a biological experiment

$$H_{ex} : \{a = 0, b = 1, c = 0\} \quad b- ; a+ ; c+ ; a- ; b+ \quad \{a = 0, b = 1, c = 1\}$$

(A genetically modified Hoare logic,
Bernot et al., 2019)

↓

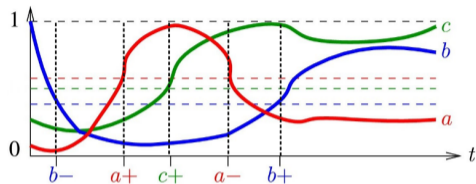
New postcondition noted Q_1 :

$$\underline{\mathbf{K}_{b,\omega} \geq 1} \wedge a = 0 \wedge b = 0 \wedge c = 1$$

The genetically modified Hoare logic

Hoare triple noted H: {Pre} Path {Post}

- Precondition: $a=0, b=1, c=0$
- Path: $b-; a+; c+; a-; b+$
- Postcondition: $a=0, b=1, c=1$



Normalised expression profiles from a biological experiment

$$H_{ex} : \{a = 0, b = 1, c = 0\} \quad b- ; a+ ; c+ ; a- ; b+ \quad \{a = 0, b = 1, c = 1\}$$



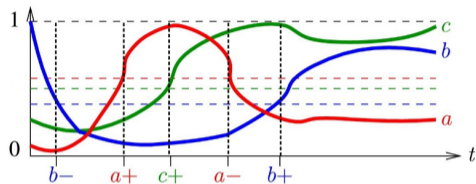
New postcondition noted Q_2 :

$$\underline{(\mathbf{K}_{b,\omega} \geq \mathbf{1}) \wedge (\mathbf{K}_{a,\omega} < \mathbf{1})} \wedge a = \mathbf{1} \wedge b = \mathbf{0} \wedge c = \mathbf{1}$$

The genetically modified Hoare logic

Hoare triple noted H: {Pre} Path {Post}

- Precondition: $a=0, b=1, c=0$
- Path: $b-; a+; c+; a-; b+$
- Postcondition: $a=0, b=1, c=1$



Normalised expression profiles from a biological experiment

$$H_{ex} : \{a = 0, b = 1, c = 0\} \quad b- ; a+ ; c+ ; a- ; b+ \quad \{a = 0, b = 1, c = 1\}$$

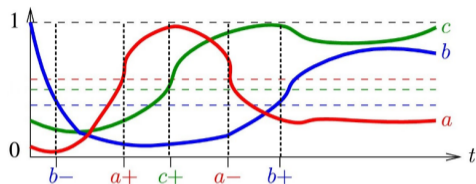

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The genetically modified Hoare logic

Hoare triple noted H: {Pre} Path {Post}

- Precondition: $a=0, b=1, c=0$
- Path: $b-; a+; c+; a-; b+$
- Postcondition: $a=0, b=1, c=1$



Normalised expression profiles from a biological experiment

$$H_{ex} : \{a = 0, b = 1, c = 0\} \quad b- ; a+ ; c+ ; a- ; b+ \quad \{a = 0, b = 1, c = 1\}$$

↓
 Q_5

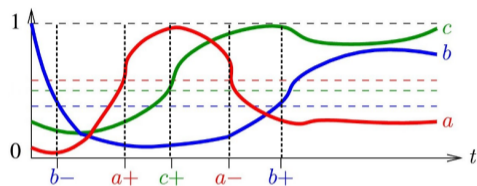
Weakest Precondition (WP)

$$\underline{(\mathbf{K}_{b,\omega} \geq 1) \wedge (\mathbf{K}_{a,\omega} < 1) \wedge (\mathbf{K}_{c,\omega} \geq 1) \wedge (\mathbf{K}_{a,\omega} \geq 1) \wedge (\mathbf{K}_{b,\omega} < 1)} \wedge a = 0 \wedge b = 1 \wedge c = 0$$

The genetically modified Hoare logic

Hoare triple noted H: $\{Pre\}$ Path $\{Post\}$

- Precondition: $a=0, b=1, c=0$
- Path: $b-; a+; c+; a-; b+$
- Postcondition: $a=0, b=1, c=1$



Normalised expression profiles from a biological experiment

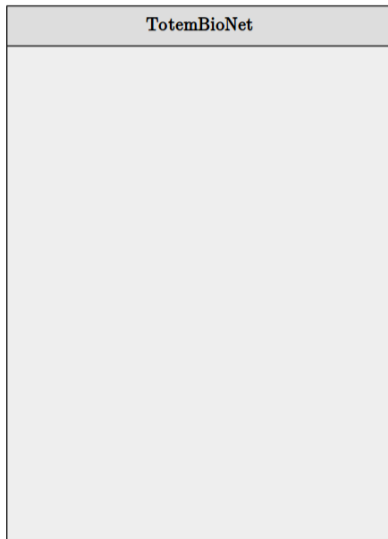
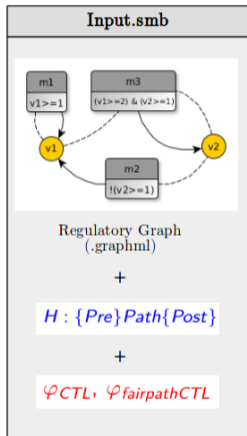
Path: $b-; a+; \exists((a+; c+), (c+; a+)); a-; b+$

↓
Disjunctive WP

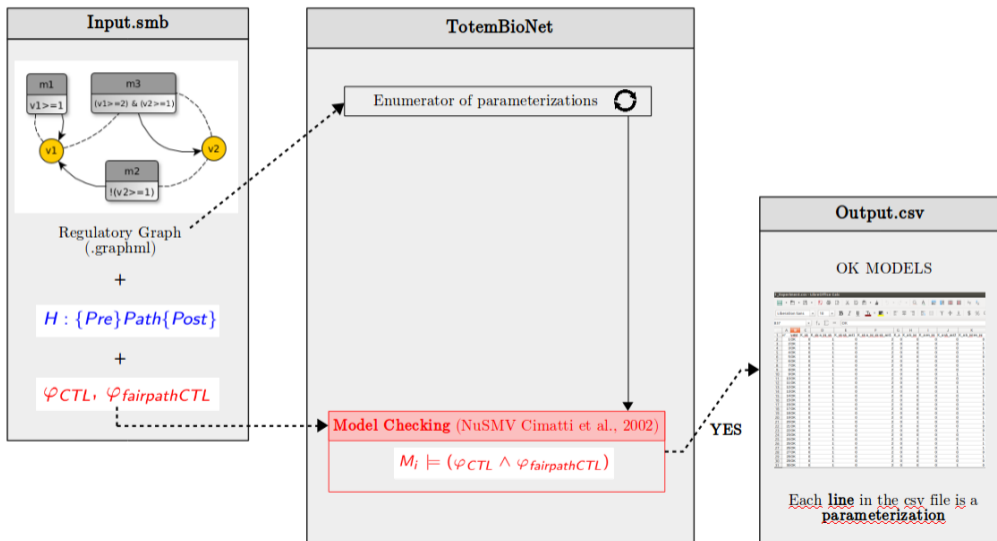
Path: $b-; a+; \forall((a+; c+), (c+; a+)); a-; b+$

↓
Conjunctive WP

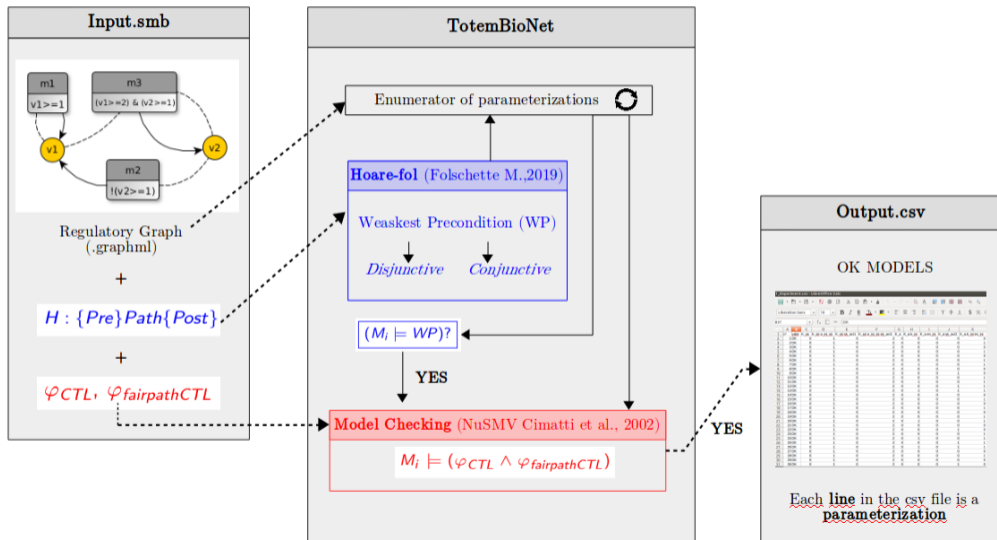
TotemBioNet workflow



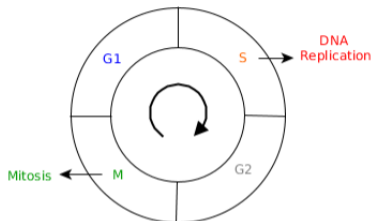
TotemBioNet workflow



TotemBioNet workflow

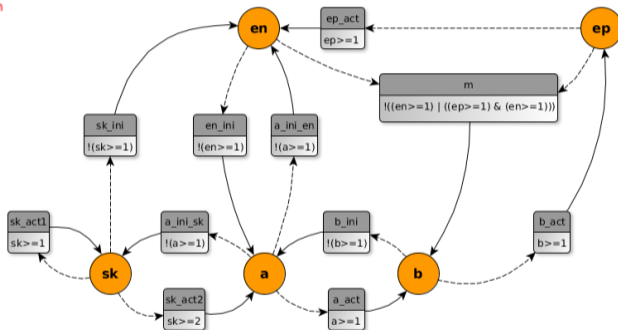


A multivalued mammalian cell cycle with 5 abstract variables



- **sk** : CycE/Cdk2
- **a** : CycA/Cdk1
- **b** : CycB/Cdk1
- **en** : APC-cdh1, Wee1, p21, p27
- **ep** : APC-cdc20

Behaegel et al., JBCB 2016



DEMO 1 DEMO 2

Graph made with yEd : <https://www.yworks.com/products/yed>

Verification of the cell cycle with TotemBioNet

Cell cycle and phases (G1,S,G2,M):

$$H_{init} : \left\{ G1_{init} \right\} \begin{array}{l} sk+; sk+; en-; \\ a+; sk-; sk-; b+; \\ a-; ep+; \\ en+; b-; ep-; \end{array} \left\{ G1_{init} \right\}$$

A globally cyclic behaviour:

$$\varphi_{cyclic} \equiv G1_{init} \Rightarrow AX(AF(G1_{init}))$$

with $G1_{init}$ the state $sk = 0, ep = 0, a = 0, b = 0, en = 1$.

Experiment	Hoare triple	$ H ^1$	Temporal logic formula	$ S ^3$	Computation Time (s) ²
1 (DEMO)	H_{init}	676	φ_{cyclic}	609	6.1

¹ $|H|$: number of parameterizations satisfying the Hoare triple, $|S|$ satisfying both Hoare triple and formulas

²Performed on an Intel Core i7-8650U processor, 1.90GHz, 8 cores.

Verification of an hypothesis about cell cycle phases with TotemBioNet

$$\mathbf{H}_{perm} : \left\{ G1_{init} \right\}$$

$Forall((sk+; sk+; en-), (sk+; en-; sk+), (en-; sk+; sk+));$
 $Forall((a+; sk-; sk-; b+), (a+; sk-; b+; sk-), (a+; b+; sk-; sk-), (sk-; a+; sk-; b+),$
 $(sk-; a+; b+; sk-), (b+; a+; sk-; sk-), (sk-; sk-; a+; b+), (sk-; b+; a+; sk-),$
 $(b+; sk-; a+; sk-), (sk-; sk-; b+; a+), (sk-; b+; sk-; a+), (b+; sk-; sk-; a+))$
 $Forall((ep+; a-), (a-; ep+));$
 $Forall((en+; b-; ep-), (en+; ep-; b-), (ep-; b-; en+),$
 $(ep-; en+; b-), (b-; en+; ep-), (b-; ep-; en+));$

$$\left\{ G1_{init} \right\}$$

Exp	Hoare triple	$ H ^3$	Temporal logic formula	$ S ^5$	Computation Time (s)
2 (DEMO)	H_{perm}	0	φ_{cyclic}	0	0.24

³ $|H|$: number of parameterizations satisfying the Hoare triple, $|S|$ satisfying both Hoare triple and formulas

Verification of an hypothesis about cell cycle phases with TotemBioNet

$$\begin{aligned}
 & \text{Forall}((sk+; sk+; en-), (sk+; en-; sk+), (en-; sk+; sk+)); \\
 & \text{Forall}((a+; sk-; sk-; b+), (a+; sk-; b+; sk-), (a+; b+; sk-; sk-), (sk-; a+; sk-; b+), \\
 & \quad (sk-; a+; b+; sk-), (b+; a+; sk-; sk-), (sk-; sk-; a+; b+), (sk-; b+; a+; sk-), \\
 & \quad (b+; sk-; a+; sk-), (sk-; sk-; b+; a+), (sk-; b+; sk-; a+), (b+; sk-; sk-; a+)) \quad \{ G1_{init} \} \\
 & \quad \text{Forall}((ep+; a-), (a-; ep+)); \\
 & \quad \text{Forall}((en+; b-; ep-), (en+; ep-; b-), (ep-; b-; en+), \\
 & \quad (ep-; en+; b-), (b-; en+; ep-), (b-; ep-; en+)); \\
 \\
 & \text{Forall}((sk+; sk+; en-), (sk+; en-; sk+), (en-; sk+; sk+)); \\
 & \text{Forall}((sk-; sk-; b+), (sk-; b+; sk-), (b+; sk-; sk-)); \\
 & \text{Forall}((ep+; a-), (a-; ep+)); \\
 & \quad en+; b-; ep-;
 \end{aligned}$$

$$\begin{aligned}
 H_{perm} : \{ G1_{init} \} & \\
 \\
 H_{permG1} : \{ G1_{init} \} &
 \end{aligned}$$

Exp	Hoare triple	$ H ^4$	Temporal logic formula	$ S ^5$	Computation Time (s)
2 (DEMO)	H_{perm}	0	φ_{cyclic}	0	0.24
3 (DEMO)	H_{permG1}	260	φ_{cyclic}	240	2.4

⁴ $|H|$: number of parameterizations satisfying the Hoare triple, $|S|$ satisfying both Hoare triple and formulas

Verification of cell cycle checkpoints with TotemBioNet

$$\mathbf{H}_{\text{permG1}} : \left\{ G1_{\text{init}} \right\} \quad \begin{array}{l} \underline{\text{Forall}}((sk+; sk+; en-), (sk+; en-; sk+), (en-; sk+; sk+)); \\ a+; \underline{\text{Forall}}((sk-; sk-; b+), (sk-; b+; sk-), (b+; sk-; sk-)); \\ \underline{\text{Forall}}((ep+; a-), (a-; ep+)); \\ en+; b-; ep-; \end{array} \quad \left\{ G1_{\text{init}} \right\}$$

$$\varphi_{G1/S} \equiv \left(G1_{\text{init}} \right) \Rightarrow \neg \left(\begin{array}{l} EX(a=1 \wedge EX(sk = 1 \wedge EX(en = 0 \wedge EX(sk = 2)))) \\ \vee EX(sk = 1 \wedge EX(a=1 \wedge EX(en = 0 \wedge EX(sk = 2)))) \\ \vee EX(sk = 1 \wedge EX(en = 0 \wedge EX(a=1 \wedge EX(sk = 2)))) \\ \vee EX(a=1 \wedge EX(en = 0 \wedge EX(sk = 1 \wedge EX(sk = 2)))) \\ \vee EX(sk = 1 \wedge EX(a=1 \wedge EX(sk = 1 \wedge EX(en = 0)))) \\ \vee EX(sk = 1 \wedge EX(sk = 2 \wedge EX(a=1 \wedge EX(en = 0)))) \\ \vee EX(a=1 \wedge EX(sk = 1 \wedge EX(sk = 2 \wedge EX(en = 0)))) \\ \vee EX(en = 0 \wedge EX(a=1 \wedge EX(sk = 1 \wedge EX(sk = 2)))) \\ \vee EX(en = 0 \wedge EX(sk = 1 \wedge EX(a=1 \wedge EX(sk = 2)))) \end{array} \right)$$

- with $G1_{\text{init}}$: $sk = 0, ep = 0, a = 0, b = 0, en = 1$.
- $EX(a=1 \wedge EX(sk = 1 \wedge EX(en = 0 \wedge EX(sk = 2))))$ is equivalent to the path $a+; sk+; en-; sk+$

Verification of cell cycle checkpoints with TotemBioNet

Exp	Hoare triple	$ H $	Temporal logic formula	$ S $	Computation Time (s)
3	H_{permG1}	260	$\varphi_{cyclic} \wedge \varphi_{G2/M} \wedge \varphi_{G1/S}$	28	2.9

TotemBioNet's features in a nutshell

WHAT: automates parameters' identification using two formal methods: **Hoare logic** and **fair path/CTL combined with model-checking**

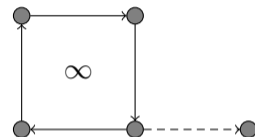
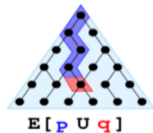
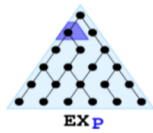
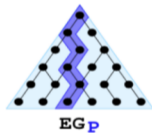
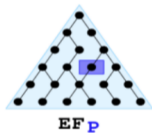
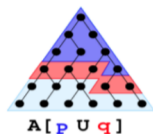
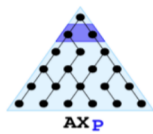
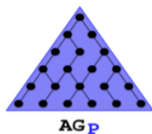
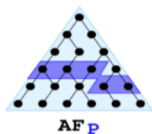
PURPOSE: formalization biological knowledge and their **quick verification**

WHERE: <https://gitlab.com/totembionet/totembionet>⁵

IMPROVEMENTS: incremental analysis of parameterizations + a Jupyter notebook.

⁵only on Linux and Mac

CTL and fair-path CTL



- p and q two properties
- Temporal modalities made up of 2 letters : a *quantifier* and a *temporal operator*
- **Quantifiers:** A,E, **Temporal operators:** F,G,X,U